

Multi-Agent Systems and Social Influence

Part 2

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(Only apparently overwhelming) plan

- 1 Fenrong Liu, Jeremy Seligman, Patrick Girard:

Logical dynamics of belief change in the community

- 2 Zoé Christoff, Jens Ulrik Hansen:

A Two-Tiered Formalization of Social Influence

- 3 Alexandru Baltag, Zoé Christoff, Rasmus Rendsvig, Sonja Smets

Dynamic Epistemic Logic of Diffusion and Prediction in Threshold Models

- 4 Truls Pedersen, Marija Slavkovik

Formal Models of Conflicting Social Influence

- 5 Sonja Smets, Fernando Velázquez-Quesada

How to Make Friends: A Logical Approach to Social Group Creation

Logical dynamics of belief change in the community

- Explores formalization of the relationship between
 - ▶ *norms of belief revision*, and
 - ▶ *properties of networks*.

Modal Logic in two minutes

- Simple language

$$\phi ::= p \mid \neg\phi \mid \phi \wedge \phi \mid \Box\phi$$

- Usually extended

$$\phi \vee \psi := \neg(\neg\phi \wedge \neg\psi)$$

$$\Diamond\phi := \neg\Box\neg\phi$$

Modal Logic in two minutes

- Models (W, R, V)
 - ▶ W non-empty set of worlds,
 - ▶ $R \subseteq W \times W$, and
 - ▶ $V : W \rightarrow \wp(P)$.

Modal Logic in two minutes

- Models (W, R, V)
 - ▶ W non-empty set of worlds,
 - ▶ $R \subseteq W \times W$, and
 - ▶ $V : W \rightarrow \wp(P)$.
- Models, *in some world*, may *satisfy* a formula

$$M, w \models p \Leftrightarrow p \in V(w)$$

$$M, w \models \neg\phi \Leftrightarrow \text{it is not the case that } M, w \models \phi$$

$$\vdots$$

$$M, w \models \Box\phi \Leftrightarrow \text{for every } v \in W \text{ s.t. } wRv : M, v \models \phi$$

Hybrid Logic in two minutes

- In addition to regular propositional symbols, NOM is a set of *names*
 - ▶ $name : NOM \rightarrow W$ names the worlds
- Behave similar to propositions, but

$$M, w \models p \Leftrightarrow p \in V(w)$$

$$M, w \models i \Leftrightarrow w = name(i)$$

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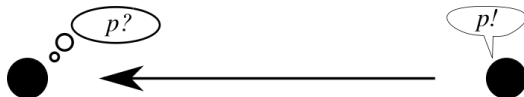
$$M, w \models i \Leftrightarrow w = name(i)$$

- Usually add *global modality* G

$$M, w \models G\phi \Leftrightarrow \text{for every } v \ M, v \models \phi$$

Doxastic Influence

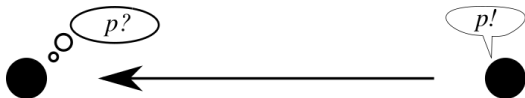
“To be influenced by my friends is to change my beliefs so that they correspond better to theirs. To begin with, we will consider influence regarding a single proposition p . If I do not believe p and some significant number or proportion of my friends do believe it, there are several ways I could respond. I could, of course, ignore their opinions and remain doxastically unperturbed. But if I am influenced to change my beliefs there are at least two ways of doing so: I may revise so that I too believe p or (more cautiously) merely contract, removing my belief in its negation $\neg p$.”



Doxastic Influence

Will I change my belief?

- 1 my own attitude regarding p ,
- 2 the cohesiveness of my *friends'* beliefs concerning p , and
- 3 the extent to which I regard any particular friend as an authority on p .



Two kinds of influence

We consider two *kinds of influence*:

Strong If we are “*strongly influenced to believe p* ”, then we *revise* in favour of p .

We denote this $S(p)$, or simply Sp .

Weak If we are “*weakly influenced to believe p* ”, then we *contract* our (possible) belief in $\neg p$.

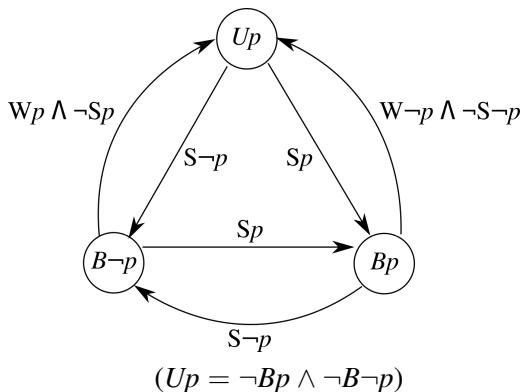
We denote this $W(p)$, or simply Wp .

Consequence of influence

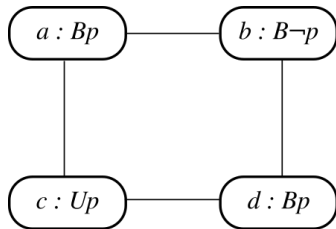
```
if  $S(p)$  then  
     $R(p)$   
else  
    if  $W(p)$  then  
         $C(\neg p)$   
    end if  
end if
```

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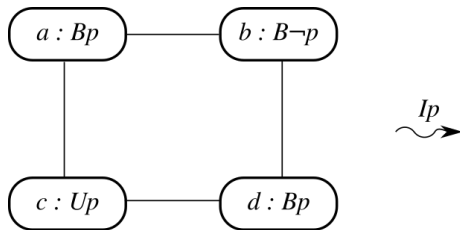


Influence as an action/event



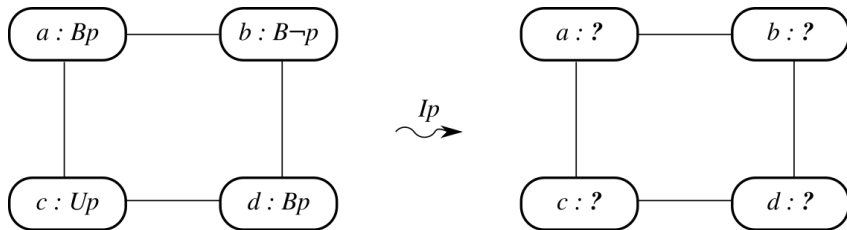
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Influence as an action/event



- We are considering how a *given network* evolves,
- an event is an operation which *updates* the model,
- want to describe the *rational outcome*.

A first instantiation

Threshold influence (conservative)

- Recall threshold models: if the fraction of your friends that believe p exceeds some threshold θ , you are influenced to believe p .

Conservative threshold model

Sp If all your friends believe p , and you have at least one friend.

Wp If at least one of your friends believe p , and none believe $\neg p$.

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- “Social hermits” are not affected by social influence.
- “Conservative” in the sense that our conditions are not “risky”. If belief spreads across network edges and *every friend* believes p , we must accept it. If there is *no support* for $\neg p$, but at least one friend believes p , we should not contradict her.

Logical language

Language

The language describing influence is given by

$$\phi ::= \underbrace{Bp \mid B\neg p \mid Sp \mid S\neg p \mid Wp \mid W\neg p}_{\text{atoms}} \mid \underbrace{\neg\phi \mid \phi \wedge \phi}_{\text{boolean}} \mid \underbrace{F\phi \mid \langle F \rangle \phi}_{\text{new}}$$

$F\phi$ For every friend, ϕ .

$\langle F \rangle \phi$ For some friend, ϕ .

($\langle F \rangle \phi := \neg F\neg\phi$ can be treated as an abbreviation.)

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Conservative threshold model formalized

$$Sp \Leftrightarrow FBp \wedge \langle F \rangle Bp$$

$$Wp \Leftrightarrow F\neg B\neg p \wedge \langle F \rangle Bp$$

Reading formulas

Assume we are discussing a particular agent, a , with “friends” $N(a)$

$$Sp \Leftrightarrow FBp \wedge \langle F \rangle Bp$$

\Leftrightarrow every $b \in N(a)$ believes p , and

there is at least one $b \in N(a)$ that believes p

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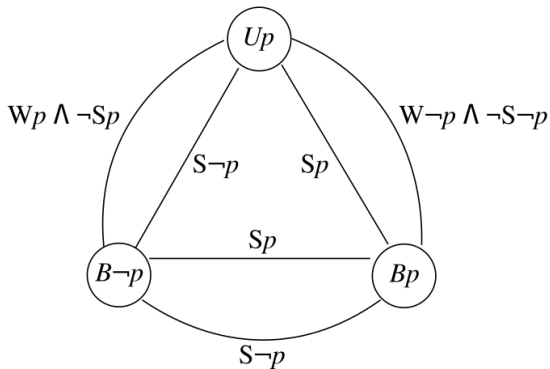
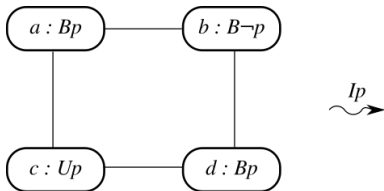
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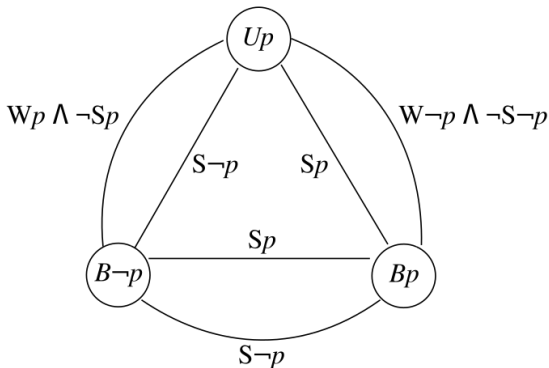
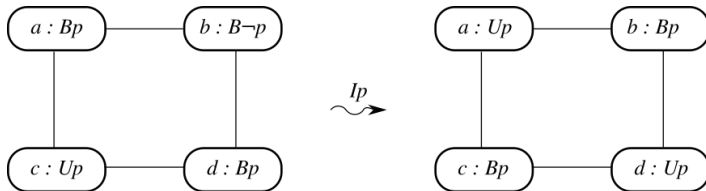
$$\Leftrightarrow F\neg B\neg p \wedge \langle F \rangle Bp$$

$$\Leftrightarrow Wp$$

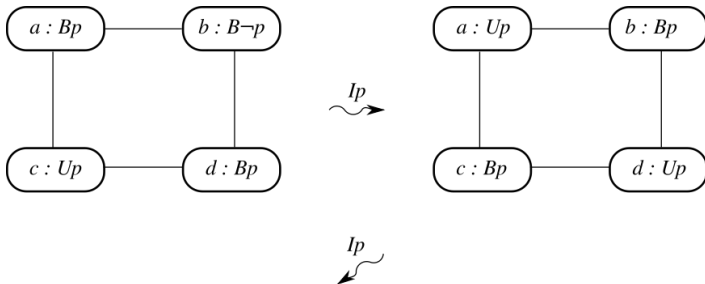
Model Updates



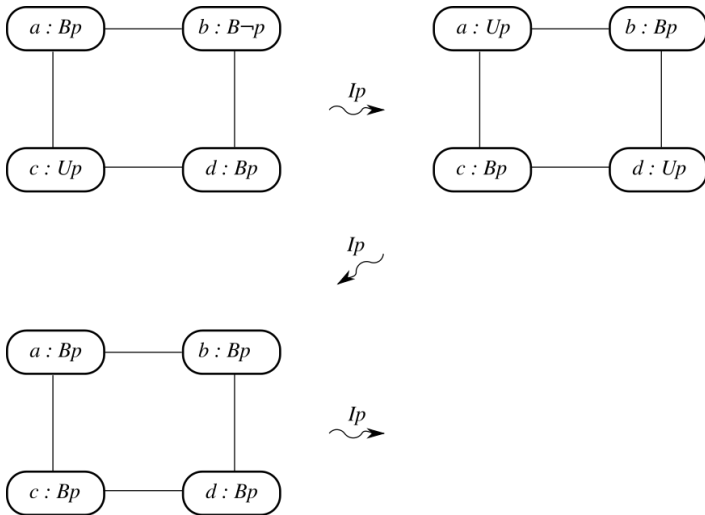
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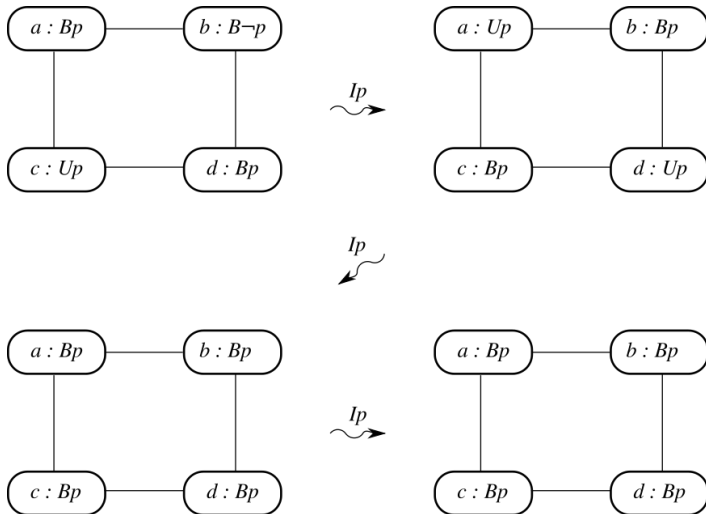
Model Updates



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Stability and flux

Stable Models that are “unaffected” by the update operation, are *stable*
They are *fixed points* of the update operation.

$$\mathcal{M} = I_p(\mathcal{M})$$

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Becoming stable Models that, after a number of applications of the update operator, become such a fixed point, are *becoming stable*.

$$I_p^n(\mathcal{M}) = I_p(I_p^n(\mathcal{M}))$$

Stability and flux

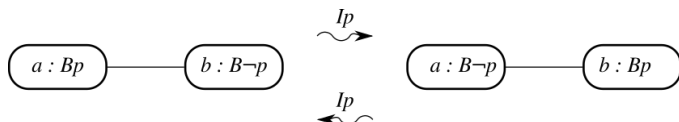
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In flux The remaining models are *in flux*.



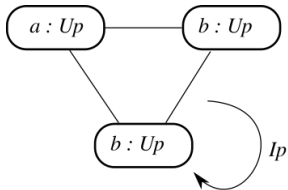
Local conditions of stability

- If no agent changes, the network is *stable*.

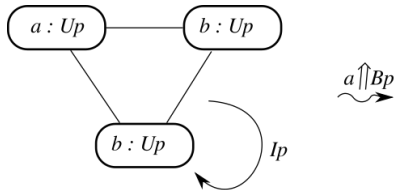
$$\neg(B\neg p \wedge Wp) \wedge \neg(Up \wedge Sp) \wedge \neg(Up \wedge S\neg p) \wedge \neg(Bp \wedge W\neg p)$$

- Network is stable if, and only if, every agent satisfies the above.
 - ▶ (We already saw $Sp \Rightarrow Wp$.)

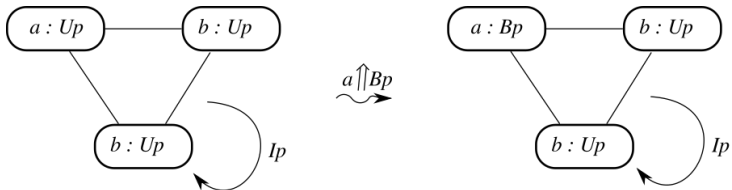
Unilateral Update



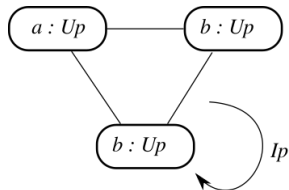
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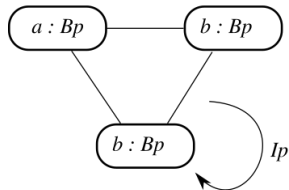
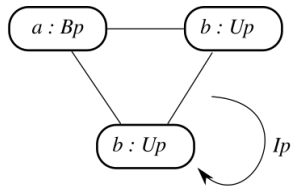
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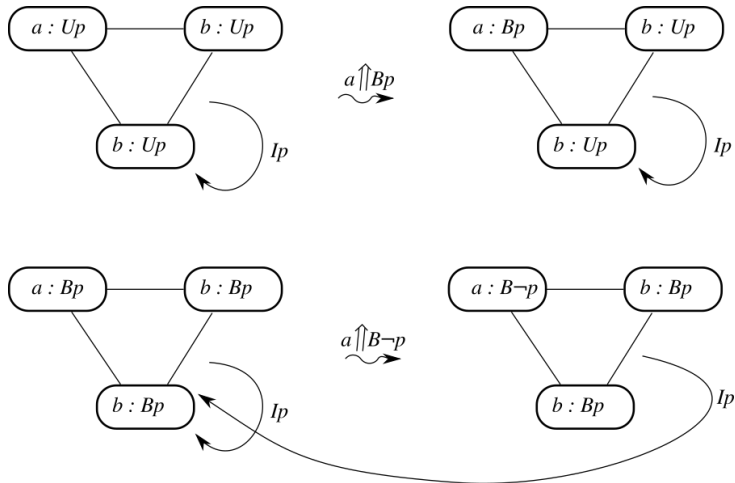


$a \uparrow \parallel Bp$

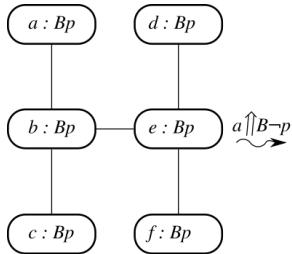


$a \uparrow \parallel B \neg p$

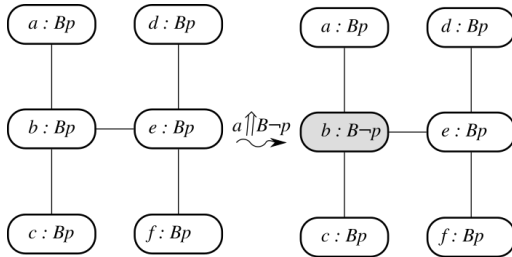
Unilateral Update



n -resistant networks

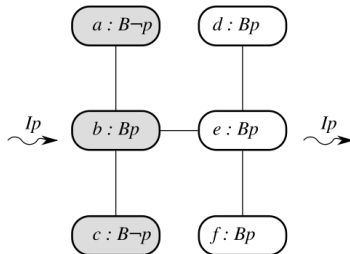
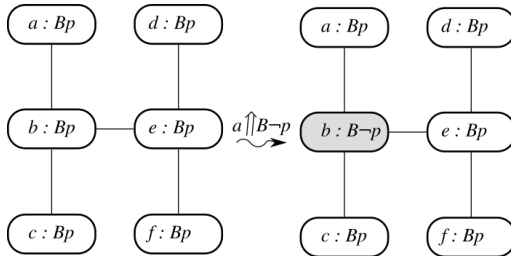


n -resistant networks

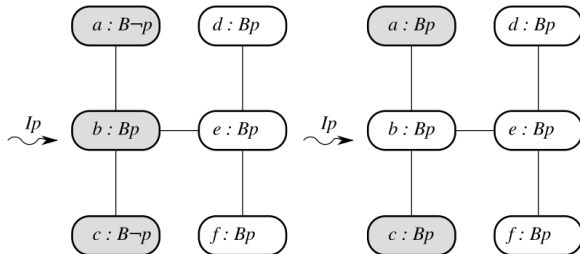
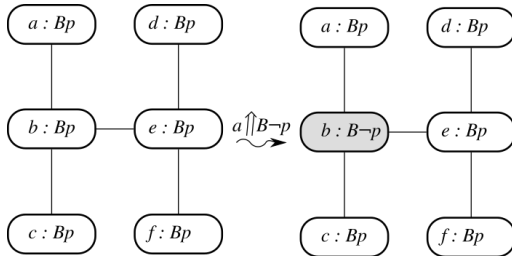


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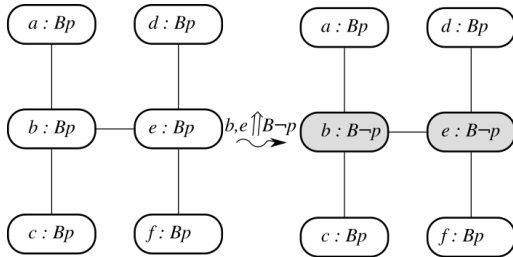
n -resistant networks



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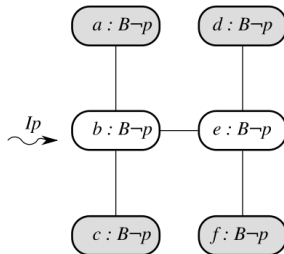
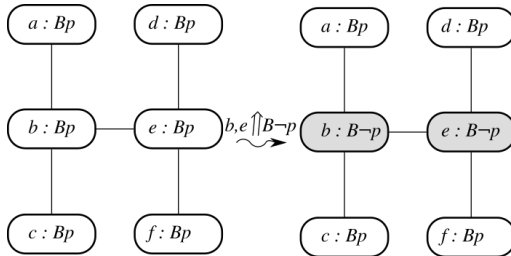


n -resistant networks



Ip

n -resistant networks



Logical dynamics of belief change in the community

Liu, Seligman & Girard

- Establishes general terminology and preliminary issues
 - ▶ strong and weak influence,
 - ▶ position in network matters,
 - ▶ n -resistance,
 - ▶ timing (consecutive \neq simultaneous),
 - ▶ conditions for stability.
- Goes on to discuss
 - ▶ Ranking friends' reliability (degree of influence),
 - ▶ possibility of *chaning* the network
 - ★ adding/removing edges
 - ▶ Appendix discusses similar investigations in other fields

A Two-Tiered Formalization of Social Influence

Zoé Christoff and Jens Ulrik Hansen

“The [framework of Liu, Seligman, and Girard] makes it unproblematic to identify the stability and stabilization conditions of social-doxastic configurations, both of which can be characterized directly in the language of friendship and belief. However, this simplicity is pricey: [...] it relies on an extremely strong assumption: agents’ belief states are influenced directly by their friends’ belief states.

Thus, either all agents have direct access to their friends’ beliefs (as mind-readers would), or their observed behavior always reflects their private beliefs, i.e., there is no difference between what they seem to believe and what they actually believe.”

Pluralistic Ignorance

Pluralistic Ignorance

- Agents believe that their private attitudes differ from the others', *but*
 - agents' private attitudes actually coincide.
-
- Need a *two-tiered* model:
 - ▶ *private (true) belief* (dislikes hats), and
 - ▶ *expressed/observable behavior* (wears a hat).
 - We expand the framework, but keep
 - ▶ three states of belief (Bp , $B\neg p$ and Up),
 - ▶ conservative threshold model for strong and weak influence.

Hybrid Network Logic

Language

Language

The language of *Hybrid Network Logic* is given by

$$\phi ::= \underbrace{V_l = r_j \mid i}_{\text{atoms}} \mid \underbrace{\neg\phi \mid \phi \wedge \phi}_{\text{boolean}} \mid \overbrace{F\phi}^{\text{modal}} \mid \underbrace{G\phi \mid @_i\phi}_{\text{hybrid}}$$

where

- V_l is the name of the j th *characteristic*,
- r_j is a possible value for the j th characteristic ($r_j \in R_j$), and
- $i \in \text{NOM}$ is a *nominal*.

Characteristic

There are a finite set of *characteristics*: V_1, V_2, \dots, V_n , each with an associated finite domain R_1, R_2, \dots, R_n (e.g., $R_1 = \{r_{1,1}, r_{1,2}, \dots, r_{1,m}\}$).

Hybrid Network Logic

Models

Hybrid Network Logic Model

A *Hybrid Network Logic Model* is a tuple $\mathcal{M} = (A, \sim, g, v)$ where

- A is a non-empty set of agents,
- $\sim \subseteq A \times A$ is a symmetric and irreflexive network relation,
- $g : \text{NOM} \rightarrow A$ names agents, and
- for every $a \in A$, $v(a)$ maps each characteristic to a possible value:
 $v(a)(V_l) \in R_l$.

Hybrid Network Logic

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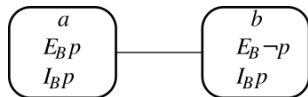
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 $v(a)(V_I) \in R_I$.

Suppose there is a characteristic V_I , with possible values $\{B_p, B_{\neg p}, U_p\}$, corresponding to the *internal belief state*. If agent a believes p , b believes $\neg p$, and agent c is undecided, we get:

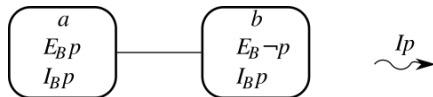
- $v(a)(V_I) = B_p$,
- $v(b)(V_I) = B_{\neg p}$, and
- $v(c)(V_I) = U_p$.

Similar possible instability



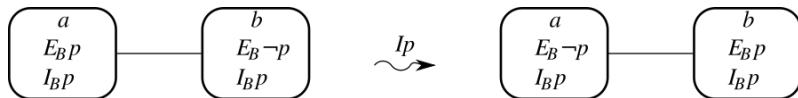
- Both agents *actually agree* (internal belief),
- each agent *apparently contradicts* the other.

Similar possible instability



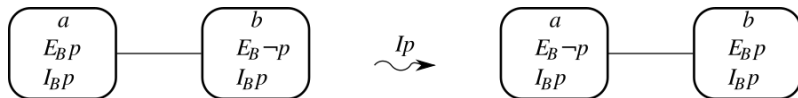
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- We transform/update the model as usual.

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- The same as before, but the other agent now expresses her inner belief.

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Now we need to decide how the *internal* beliefs affects the agents' belief revision.

Two-Tiered Norms of Behavior

| | Inner state | $\langle F \rangle E_B \varphi$ | $\langle F \rangle E_B \neg \varphi$ | $\langle F \rangle E_U \varphi$ | Type 1 | Type 2 | Type 3 |
|----|--------------------|---------------------------------|--------------------------------------|---------------------------------|-----------------------------|-----------------------------|-----------------------------|
| 1 | $I_B \varphi$ | | | | $\leadsto E_B \varphi$ | $\leadsto E_U \varphi$ | $\leadsto E_B \varphi$ |
| 2 | $I_B \neg \varphi$ | 1 | 1 | 1 | $\leadsto E_B \neg \varphi$ | $\leadsto E_U \varphi$ | $\leadsto E_B \neg \varphi$ |
| 3 | $I_U \varphi$ | | | | $\leadsto E_U \varphi$ | $\leadsto E_U \varphi$ | $\leadsto E_U \varphi$ |
| 4 | $I_B \varphi$ | | | | $\leadsto E_B \varphi$ | $\leadsto E_U \varphi$ | $\leadsto E_B \varphi$ |
| 5 | $I_B \neg \varphi$ | 1 | 1 | 0 | $\leadsto E_B \neg \varphi$ | $\leadsto E_U \varphi$ | $\leadsto E_B \neg \varphi$ |
| 6 | $I_U \varphi$ | | | | $\leadsto E_U \varphi$ | $\leadsto E_U \varphi$ | $\leadsto E_U \varphi$ |
| 7 | $I_B \varphi$ | | | | $\leadsto E_B \varphi$ | $\leadsto E_B \varphi$ | $\leadsto E_B \varphi$ |
| 8 | $I_B \neg \varphi$ | 1 | 0 | 1 | $\leadsto E_U \varphi$ | $\leadsto E_U \varphi$ | $\leadsto E_U \varphi$ |
| 9 | $I_U \varphi$ | | | | $\leadsto E_U \varphi$ | $\leadsto E_U \varphi$ | $\leadsto E_U \varphi$ |
| 10 | $I_B \varphi$ | | | | $\leadsto E_B \varphi$ | $\leadsto E_B \varphi$ | $\leadsto E_B \varphi$ |
| 11 | $I_B \neg \varphi$ | 1 | 0 | 0 | $\leadsto E_B \varphi$ | $\leadsto E_B \varphi$ | $\leadsto E_B \varphi$ |
| 12 | $I_U \varphi$ | | | | | | |
| 13 | $I_B \varphi$ | | | | $\leadsto E_U \varphi$ | $\leadsto E_U \varphi$ | $\leadsto E_U \varphi$ |
| 14 | $I_B \neg \varphi$ | 0 | 1 | 1 | $\leadsto E_B \neg \varphi$ | $\leadsto E_B \neg \varphi$ | $\leadsto E_B \neg \varphi$ |
| 15 | $I_U \varphi$ | | | | $\leadsto E_B \neg \varphi$ | $\leadsto E_B \neg \varphi$ | $\leadsto E_B \neg \varphi$ |
| 16 | $I_B \varphi$ | | | | $\leadsto E_B \neg \varphi$ | $\leadsto E_B \neg \varphi$ | $\leadsto E_B \neg \varphi$ |
| 17 | $I_B \neg \varphi$ | 0 | 1 | 0 | $\leadsto E_B \neg \varphi$ | $\leadsto E_B \neg \varphi$ | $\leadsto E_B \neg \varphi$ |
| 18 | $I_U \varphi$ | | | | | | |
| 19 | $I_B \varphi$ | | | | $\leadsto E_B \varphi$ | $\leadsto E_B \varphi$ | $\leadsto E_U \varphi$ |
| 20 | $I_B \neg \varphi$ | 0 | 0 | 1 | $\leadsto E_B \neg \varphi$ | $\leadsto E_B \neg \varphi$ | $\leadsto E_U \varphi$ |
| 21 | $I_U \varphi$ | | | | $\leadsto E_U \varphi$ | $\leadsto E_U \varphi$ | $\leadsto E_U \varphi$ |
| 22 | $I_B \varphi$ | | | | $\leadsto E_B \varphi$ | $\leadsto E_B \varphi$ | $\leadsto E_B \varphi$ |
| 23 | $I_B \neg \varphi$ | 0 | 0 | 0 | $\leadsto E_B \neg \varphi$ | $\leadsto E_B \neg \varphi$ | $\leadsto E_B \neg \varphi$ |
| 24 | $I_U \varphi$ | | | | $\leadsto E_U \varphi$ | $\leadsto E_U \varphi$ | $\leadsto E_U \varphi$ |

Two-Tiered Norms of Behavior

| | Inner state | $\langle F \rangle E_B \varphi$ | $\langle F \rangle E_B \neg \varphi$ | $\langle F \rangle E_U \varphi$ | Type 1 | Type 2 | Type 3 |
|----|--------------------|---------------------------------|--------------------------------------|---------------------------------|-----------------------------|-----------------------------|-----------------------------|
| 1 | $I_B \varphi$ | | | | $\leadsto E_B \varphi$ | $\leadsto E_U \varphi$ | $\leadsto E_B \varphi$ |
| 2 | $I_B \neg \varphi$ | 1 | 1 | 1 | $\leadsto E_B \neg \varphi$ | $\leadsto E_U \varphi$ | $\leadsto E_B \neg \varphi$ |
| 3 | $I_U \varphi$ | | | | $\leadsto E_U \varphi$ | $\leadsto E_U \varphi$ | $\leadsto E_U \varphi$ |
| 4 | $I_B \varphi$ | | | | $\leadsto E_B \varphi$ | $\leadsto E_U \varphi$ | $\leadsto E_B \varphi$ |
| 5 | $I_B \neg \varphi$ | 1 | 1 | 0 | $\leadsto E_B \neg \varphi$ | $\leadsto E_U \varphi$ | $\leadsto E_B \neg \varphi$ |
| 6 | $I_U \varphi$ | | | | $\leadsto E_U \varphi$ | $\leadsto E_U \varphi$ | $\leadsto E_U \varphi$ |
| 7 | $I_B \varphi$ | | | | $\leadsto E_B \varphi$ | $\leadsto E_B \varphi$ | $\leadsto E_B \varphi$ |
| 8 | $I_B \neg \varphi$ | 1 | 0 | 1 | $\leadsto E_U \varphi$ | $\leadsto E_U \varphi$ | $\leadsto E_U \varphi$ |
| 9 | $I_U \varphi$ | | | | $\leadsto E_U \varphi$ | $\leadsto E_U \varphi$ | $\leadsto E_U \varphi$ |
| 10 | $I_B \varphi$ | | | | $\leadsto E_B \varphi$ | $\leadsto E_B \varphi$ | $\leadsto E_B \varphi$ |
| 11 | $I_B \neg \varphi$ | 1 | 0 | 0 | $\leadsto E_B \varphi$ | $\leadsto E_B \varphi$ | $\leadsto E_B \varphi$ |
| 12 | $I_U \varphi$ | | | | $\leadsto E_B \varphi$ | $\leadsto E_B \varphi$ | $\leadsto E_B \varphi$ |
| 13 | $I_B \varphi$ | | | | $\leadsto E_U \varphi$ | $\leadsto E_U \varphi$ | $\leadsto E_U \varphi$ |
| 14 | $I_B \neg \varphi$ | 0 | 1 | 1 | $\leadsto E_B \neg \varphi$ | $\leadsto E_B \neg \varphi$ | $\leadsto E_B \neg \varphi$ |
| 15 | $I_U \varphi$ | | | | $\leadsto E_B \neg \varphi$ | $\leadsto E_B \neg \varphi$ | $\leadsto E_B \neg \varphi$ |
| 16 | $I_B \varphi$ | | | | $\leadsto E_B \neg \varphi$ | $\leadsto E_B \neg \varphi$ | $\leadsto E_B \neg \varphi$ |
| 17 | $I_B \neg \varphi$ | 0 | 1 | 0 | $\leadsto E_B \neg \varphi$ | $\leadsto E_B \neg \varphi$ | $\leadsto E_B \neg \varphi$ |
| 18 | $I_U \varphi$ | | | | $\leadsto E_B \neg \varphi$ | $\leadsto E_B \neg \varphi$ | $\leadsto E_B \neg \varphi$ |
| 19 | $I_B \varphi$ | | | | $\leadsto E_B \varphi$ | $\leadsto E_B \varphi$ | $\leadsto E_U \varphi$ |
| 20 | $I_B \neg \varphi$ | 0 | 0 | 1 | $\leadsto E_B \neg \varphi$ | $\leadsto E_B \neg \varphi$ | $\leadsto E_U \varphi$ |
| 21 | $I_U \varphi$ | | | | $\leadsto E_U \varphi$ | $\leadsto E_U \varphi$ | $\leadsto E_U \varphi$ |
| 22 | $I_B \varphi$ | | | | $\leadsto E_B \varphi$ | $\leadsto E_B \varphi$ | $\leadsto E_B \varphi$ |
| 23 | $I_B \neg \varphi$ | 0 | 0 | 0 | $\leadsto E_B \neg \varphi$ | $\leadsto E_B \neg \varphi$ | $\leadsto E_B \neg \varphi$ |
| 24 | $I_U \varphi$ | | | | $\leadsto E_U \varphi$ | $\leadsto E_U \varphi$ | $\leadsto E_U \varphi$ |

Two-Tiered Norms of Behavior

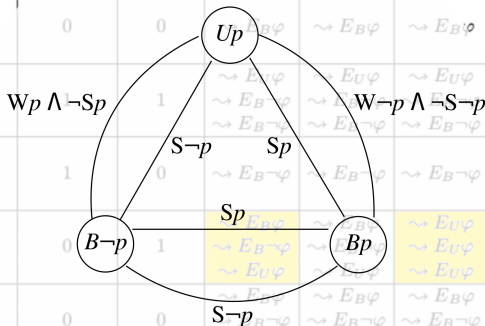
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| 4 | $I_B \varphi$ | | | | $\leadsto E_B \varphi$ | $\leadsto E_U \varphi$ | $\leadsto E_B \varphi$ |
| 5 | $I_B \neg \varphi$ | 1 | 1 | 0 | $\leadsto E_B \neg \varphi$ | $\leadsto E_U \varphi$ | $\leadsto E_B \neg \varphi$ |
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| 9 | $I_U \varphi$ | | | | $\leadsto E_U \varphi$ | $\leadsto E_U \varphi$ | $\leadsto E_U \varphi$ |
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| 16 | $I_B \varphi$ | | | | | | |
| 17 | $I_B \neg \varphi$ | 0 | 1 | 0 | $\leadsto E_B \neg \varphi$ | $\leadsto E_B \neg \varphi$ | $\leadsto E_B \neg \varphi$ |
| 18 | $I_U \varphi$ | | | | | | |
| 19 | $I_B \varphi$ | | | | $\leadsto E_B \varphi$ | $\leadsto E_B \varphi$ | $\leadsto E_U \varphi$ |
| 20 | $I_B \neg \varphi$ | 0 | 0 | 1 | $\leadsto E_B \neg \varphi$ | $\leadsto E_B \neg \varphi$ | $\leadsto E_U \varphi$ |
| 21 | $I_U \varphi$ | | | | $\leadsto E_U \varphi$ | $\leadsto E_U \varphi$ | $\leadsto E_U \varphi$ |
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| 1 | $I_B \varphi$ | | | | $\leadsto E_B \varphi$ | $\leadsto E_U \varphi$ | $\leadsto E_B \varphi$ |
| 2 | $I_B \neg \varphi$ | 1 | 1 | 1 | $\leadsto E_B \neg \varphi$ | $\leadsto E_U \varphi$ | $\leadsto E_B \neg \varphi$ |
| 3 | $I_U \varphi$ | | | | $\leadsto E_U \varphi$ | $\leadsto E_U \varphi$ | $\leadsto E_U \varphi$ |
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| 7 | $I_B \varphi$ | | | | $\leadsto E_B \varphi$ | $\leadsto E_B \varphi$ | $\leadsto E_B \varphi$ |
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| 24 | $I_U \varphi$ | | | | $\leadsto E_U \varphi$ | $\leadsto E_U \varphi$ | $\leadsto E_U \varphi$ |

Two-Tier Norms of Behavior

| | Inner state | $\langle F \rangle E_B \varphi$ | $\langle F \rangle E_B \neg \varphi$ | $\langle F \rangle E_U \varphi$ | Type 1 | Type 2 | Type 3 |
|----|--------------------|---------------------------------|--------------------------------------|---------------------------------|-------------------------|-------------------------|-------------------------|
| 1 | $I_B \varphi$ | | | | $\neg E_B \varphi$ | $\neg E_U \varphi$ | $\neg E_B \varphi$ |
| 2 | $I_B \neg \varphi$ | 1 | 1 | 1 | $\neg E_B \neg \varphi$ | $\neg E_U \varphi$ | $\neg E_B \neg \varphi$ |
| 3 | $I_U \varphi$ | | | | $\neg E_U \varphi$ | $\neg E_U \varphi$ | $\neg E_U \varphi$ |
| 4 | $I_B \varphi$ | | | | $\neg E_B \varphi$ | $\neg E_U \varphi$ | $\neg E_B \varphi$ |
| 5 | $I_B \neg \varphi$ | 1 | 1 | 0 | $\neg E_B \neg \varphi$ | $\neg E_U \varphi$ | $\neg E_B \neg \varphi$ |
| 6 | $I_U \varphi$ | | | | $\neg E_U \varphi$ | $\neg E_U \varphi$ | $\neg E_U \varphi$ |
| 7 | $I_B \varphi$ | | | | $\neg E_B \varphi$ | $\neg E_B \varphi$ | $\neg E_B \varphi$ |
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| 13 | $I_B \varphi$ | | | | $\neg E_U \varphi$ | $\neg E_U \varphi$ | $\neg E_U \varphi$ |
| 14 | $I_B \neg \varphi$ | 0 | $W_p \wedge \neg S_p$ | 1 | $\neg E_B \neg \varphi$ | $\neg E_B \neg \varphi$ | $\neg E_U \varphi$ |
| 15 | $I_U \varphi$ | | | | $\neg E_B \neg \varphi$ | $\neg E_B \neg \varphi$ | $\neg E_B \neg \varphi$ |
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| 17 | $I_B \neg \varphi$ | 0 | 1 | 0 | $\neg E_B \neg \varphi$ | $\neg E_B \neg \varphi$ | $\neg E_B \neg \varphi$ |
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New Models, New Phenomena

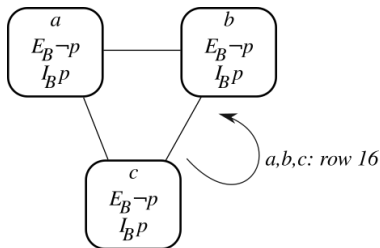
- We are now able to model *Pluralistic Ignorance*

$$PI_{\phi} \quad := \quad G(I_B\phi \wedge E_B\neg\phi)$$

New Models, New Phenomena

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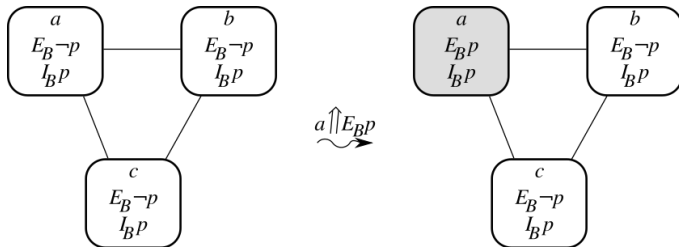


Two-Tired update

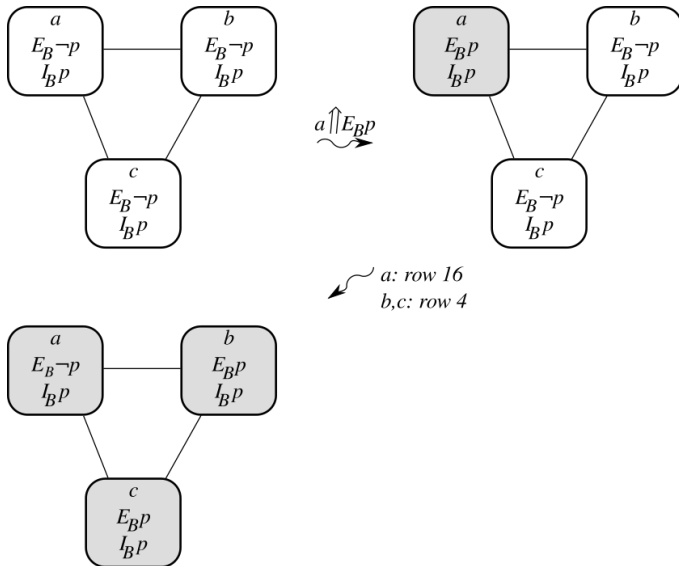
Notice rows: 4, 10, and 16.

| | Inner state | $\langle F \rangle E_B \varphi$ | $\langle F \rangle E_B \neg \varphi$ | $\langle F \rangle E_U \varphi$ | Type 1 | Type 2 | Type 3 |
|----|--------------------|---------------------------------|--------------------------------------|---------------------------------|-----------------------------|-----------------------------|-----------------------------|
| 1 | $I_B \varphi$ | | | | $\leadsto E_B \varphi$ | $\leadsto E_U \varphi$ | $\leadsto E_B \varphi$ |
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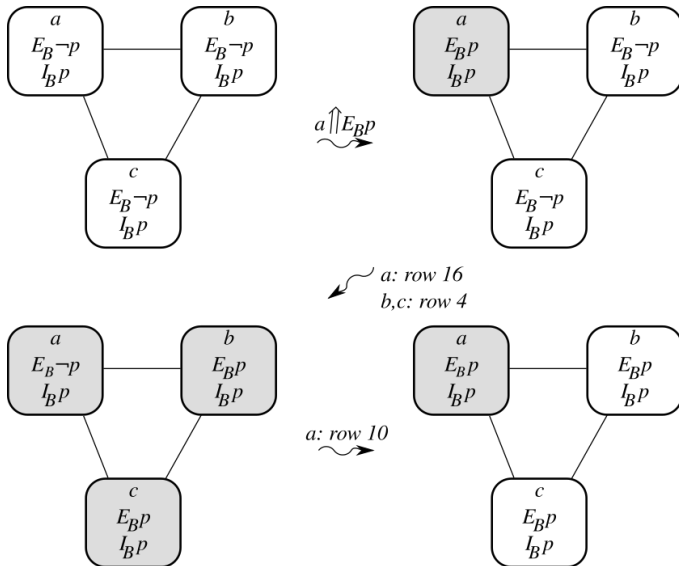
“Fragile stability”



“Fragile sability”



“Fragile sability”



A Two-Tiered Formalization of Social Influence

Christoff & Hansen

- Establishes a two-tiered hybrid logic where
 - ▶ privately held beliefs affect transition,
- demonstrates conditions for stability and (eventual) dissolution of state of pluralistic ignorance
- various *types* of agents

Dynamic Epistemic Logics of Diffusion and Prediction in Social Networks

Alexandru Baltag, Zoé Christoff, Rasmus K. Rendsvig, and Sonja Smets

“We introduce an epistemic dimension to threshld models, thus taking into account the real-life limitations posed by the agents’ limited access to information; for this, we propose an epistemic variant of the [belief/behavior] adoption rule: agents adopt a behavior only if they know that enough of [...] their neighbours have adopted it.”

Epistemic Threshold Models

- Seen simple (extreme) threshold models expressed in terms of:
 - ▶ *all*, *none*, or *some* (at least one)
- Fix a single belief/behavior we are interested in, *having adopted* p , say.
 - ▶ Denote the set agents that have adopted p by $B \subseteq \mathcal{A}$.

Epistemic Threshold Models

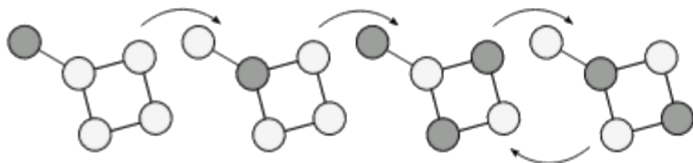
- Seen simple (extreme) threshold models expressed in terms of:
 - ▶ *all*, *none*, or *some* (at least one)
- Fix a single belief/behavior we are interested in, *having adopted* p , say.
 - ▶ Denote the set agents that have adopted p by $B \subseteq \mathcal{A}$.
- Let $\theta \in [0, 1]$ be some threshold.

$$B' = \left\{ a \in \mathcal{A} \mid \frac{|N(a) \cap B|}{|N(a)|} \geq \theta \right\}$$

- ▶ B represents initial believers,
- ▶ B' represents the result of the model update.

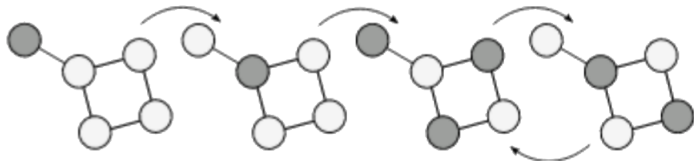
Inflationary adoption

Non-inflationary Let $\theta = 1/4$.



Inflationary adoption

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Inflationary Let $\theta = 1/4$.



Inflationary adoption

Non-inflationary

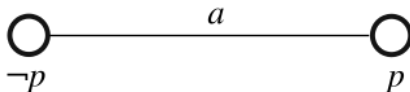
$$B' = \left\{ a \in \mathcal{A} \mid \frac{|N(a) \cap B|}{|N(a)|} \geq \theta \right\}$$

Inflationary

$$B' = B \cup \left\{ a \in \mathcal{A} \mid \frac{|N(a) \cap B|}{|N(a)|} \geq \theta \right\}$$

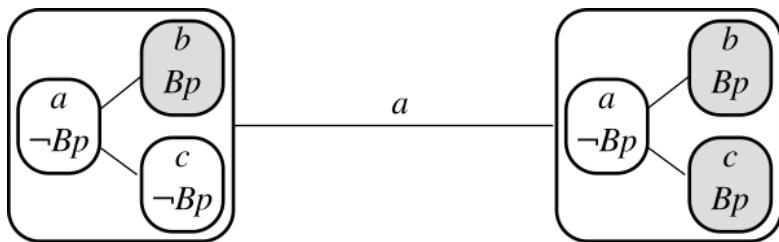
Epistemic dimension

- When should an agent adopt a belief?
 - ▶ Proposal: when she *knows* that at least θ of her friends already have adopted.



Epistemic dimension

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Epistemic Threshold Model

Epistemic Threshold Model

An ETM is a tuple $\mathcal{M} = (W, \mathcal{A}, N, B, \theta, \{\sim_a\}_{a \in \mathcal{A}})$ where

- W is a finite, non-empty set of possible worlds,
- $\sim_a \subseteq W \times W$ is an *indistinguishably* (equivalence) relation,
- $N : W \rightarrow \wp(\mathcal{A})$,
 - ▶ $a \notin N(w)(a)$ (irreflexive),
 - ▶ $b \in N(w)(a) \Leftrightarrow a \in N(w)(b)$ (symmetric), and
 - ▶ $N(w)(a) \neq \emptyset$ (serial).
- $B : W \rightarrow \wp(\mathcal{A})$

n -sight

- These models permit that agents do not “know” their friends.
(▶ like the previous example)
- If a model has 1-sight, then every agent *knows* whether every (immediate) friend has adopted p (and who they are).
 - ▶ If a model has 2-sight, then every agent *knows* whether every friend, and friends of her friends, has adopted p .
- How much knowledge does the agent need to comply to non-epistemic threshold models?

1-sight is sufficient

- *Informed update* only when the agent *knows* (*de dicto*) that the threshold is exceeded.

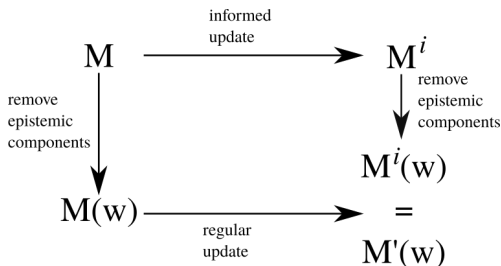
$$B'(w) = B(w) \cup \left\{ a \in \mathcal{A} \mid \forall v \sim_a w \frac{|N(v)(a) \cap B(v)|}{|N(v)(a)|} \geq \theta \right\}$$

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- *Informed update* only when the agent *knows* (*de dicto*) that the threshold is exceeded.

$$B'(w) = B(w) \cup \left\{ a \in \mathcal{A} \mid \forall v \sim_a w \frac{|N(v)(a) \cap B(v)|}{|N(v)(a)|} \geq \theta \right\}$$

- When a model has 1-sight, the following holds:



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Alexandru Baltag, Zoé Christoff, Rasmus K. Rendsvig, and Sonja Smets

- Goes on to show:
 - ▶ adding epistemic dimension can *slow down* diffusion,
 - ▶ this slowdown can be counteracted by a *speed up* when agents are permitted to *predict the development*
- provides logics for reasoning about the systems

Formal Models of Conflicting Social Influence

Truls Pedersen and Marija Slavkovik

“Assume that you have a group of friends that are convinced that climate change is a hoax. You also have another group of friends that are climate change researchers devoted to slowing down climate change. You would be under pressure to choose an opinion to support. To avoid the conflict you would necessarily have to stop your relations with at least one of (or parts of one of) the groups.”

Social Network

Social Network

A model is a tuple $\mathcal{M} = (\mathcal{A}, N, \mathcal{I}, pro)$ where

- \mathcal{I} is a set of issues (comes in pairs), e.g.,

$$\mathcal{I} = \{p, \neg p\}$$

- $pro : \mathcal{I} \rightarrow \wp(\mathcal{A})$ s.t., $pro(\phi) \cap pro(\neg\phi) = \emptyset$

- We keep three doxastic agent states: can be in either or none.

Strength of social influence

- Keep the normative question from Liu, Seligman & Girard:
when does an agent feel compelled to revise her belief?
- Primarily interested in expressing strength/degree:
by what unit do we measure this degree?
- We propose a framework for reformulating the well-known models:

$$\Omega^i : \mathcal{I} \rightarrow \wp(\wp(\mathcal{A}))$$

Pivotal sets

Pivotal sets

$A \in \Omega(i, \phi)$ is a set of i 's neighbours such that, *ceteris paribus*, after agent i drops all ties to the agents in A , i is no longer under pressure to adopt ϕ .

Pivotal sets

Pivotal sets

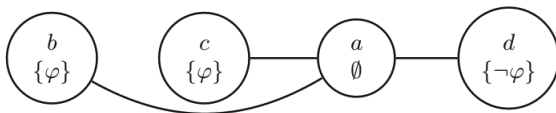
$A \in \Omega(i, \phi)$ is a set of i 's neighbours such that, *ceteris paribus*, after agent i drops all ties to the agents in A , i is no longer under pressure to adopt ϕ .

- Formalize (non-inflationary) threshold model:

$$\Omega_t(i, \phi) := \left\{ A \subseteq N(i) \mid \frac{|(N(i) \cap \text{pro}(\phi)) \setminus A|}{|N(I) \setminus A|} \not\geq \theta \right\}$$

Conflicting Social Influence

Threshold $\theta = 1/3$

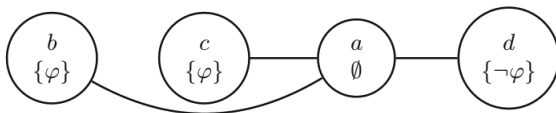


$$\Omega(a, \neg\phi) = \{\{d\}, \{b, d\}, \{c, d\}, \underbrace{\{b, c, d\}}_{N(a)}\}$$

- Can not avoid influence unless ties to d are cut.

Conflicting Social Influence

Threshold $\theta = 1/3$



$$\Omega(a, \neg\phi) = \{\{d\}, \{b, d\}, \{c, d\}, \underbrace{\{b, c, d\}}_{N(a)}\}$$

- Can not avoid influence unless ties to d are cut.

$$\Omega(a, \phi) = \{\{b, c\}, \underbrace{\{b, c, d\}}_{N(a)}\}$$

- Can not avoid influence unless ties to both b and c are cut.

Properties of Influence Models

- In both cases, $N(i)$ were included:
 - ▶ Social hermits are immune in the threshold model.
- In neither case was \emptyset included:
 - ▶ a was influenced to adopt ϕ , and influenced to adopt $\neg\phi$.
 - ▶ Provides condition for restoring dichotomous models.

Formal Models of Conflicting Social Influence

Truls Pedersen and Marija Slavkovik

- Goes on to discuss
 - ▶ temporal consequences, and
 - ▶ consistency of agents' beliefs.
- Under what conditions do elements of $\Omega(i, \phi)$ function as “conflict resolving actions”?

How to Make Friends:

A Logical Approach to Social Group Creation

Sonja Smets and Fernando R. Velázquez-Quesada

“It is commonly accepted that our social contacts affect the way we form our opinions about the world. [...] This paper focuses on the logical structure behind the creation of social networks. Our basic mechanism for group-creation focusses on agents who become socially connected when the number of features in which they differ is small enough In line with this idea we propose several group-creation policies, exploring the properties of the resulting networks.”

Making Friends

- “Birds of a feather flock together”
 - ▶ Agents that are similar may form new edges in the network.

Mismatch

The feature mismatch between a and b

$$\text{mismatch}^M(a, b) := (V(a) \setminus V(b)) \cup (V(b) \setminus V(a))$$

The mismatch *distance* between a and b

$$\text{dist}^M = |\text{mismatch}^M(a, b)|$$

Social Network Models

Social Network Models

- $\mathcal{M} = (\mathcal{A}, S, V)$
 - ▶ $S \subseteq \mathcal{A} \times \mathcal{A}$ is *any* binary relation and
 - ▶ $V : \mathcal{A} \rightarrow \wp(P)$
- Some finite P of *agent features* (propositions)
- Some updates we discuss violate the regular assumptions of symmetry and irreflexivity.

Similarity Update (I)

- We do *not* update agent features, *only* the network.

Similarity Update (I)

Given a social network (\mathcal{A}, S, V)

- S is any set of (social) edges,
- $V : \mathcal{A} \rightarrow \wp(P)$ maps every agent to the set of her features.

The similarity updated network (\mathcal{A}, S', V) with threshold θ where

$$S' = \{(a, b) \in \mathcal{A} \times \mathcal{A} \mid \text{dist}^M(a, b) \leq \theta\}$$

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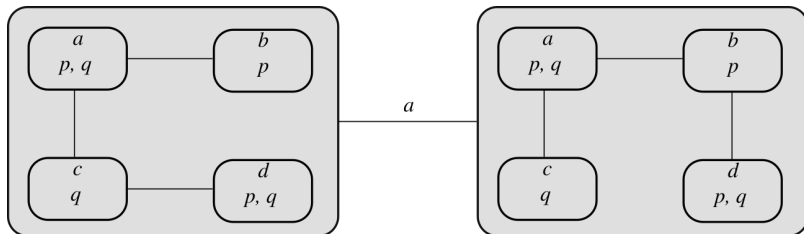
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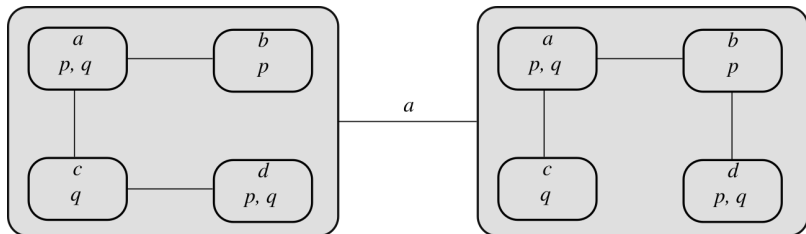


Middleman Knowledge-Based Similarity Update

Middleman Knowledge-Based Similarity Update



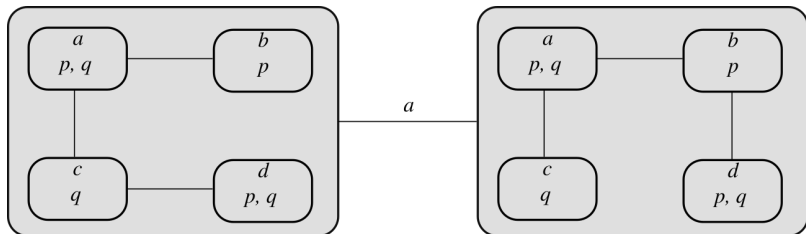
Middleman Knowledge-Based Similarity Update



de re (a, b) is added in w if there is a c such that, for every $u \sim_a w$:

- $dist_u^M(a, b)$, and
- $(a, c) \in S, (c, b) \in S$

Middleman Knowledge-Based Similarity Update



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How to Make Friends:

A Logical Approach to Social Group Creation

Sonja Smets and Fernando R. Velázquez-Quesada

- Describes several variations
 - ▶ possibility of other distances
 - ▶ with/without middleman
 - ▶ with/without epistemic dimension
- provides logics describing networks and dynamics

Questions?

Thank you!

by the way...

FYI: EUMAS2018

<https://eumas2018.w.uib.no/>

